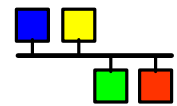


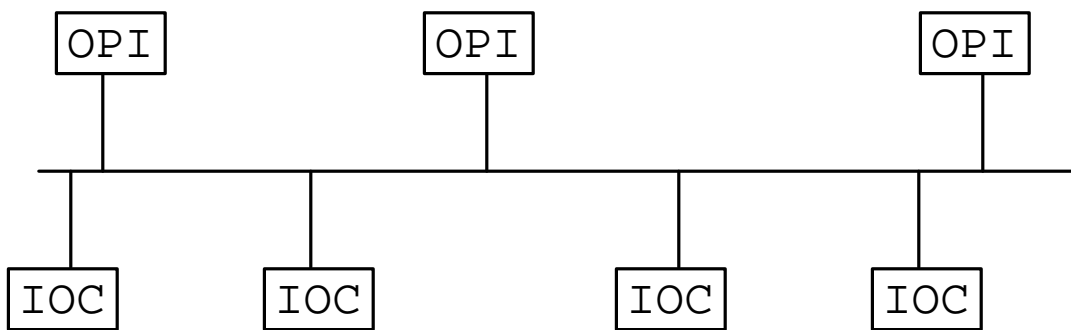
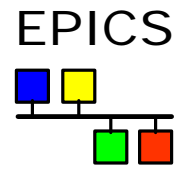
EPICS



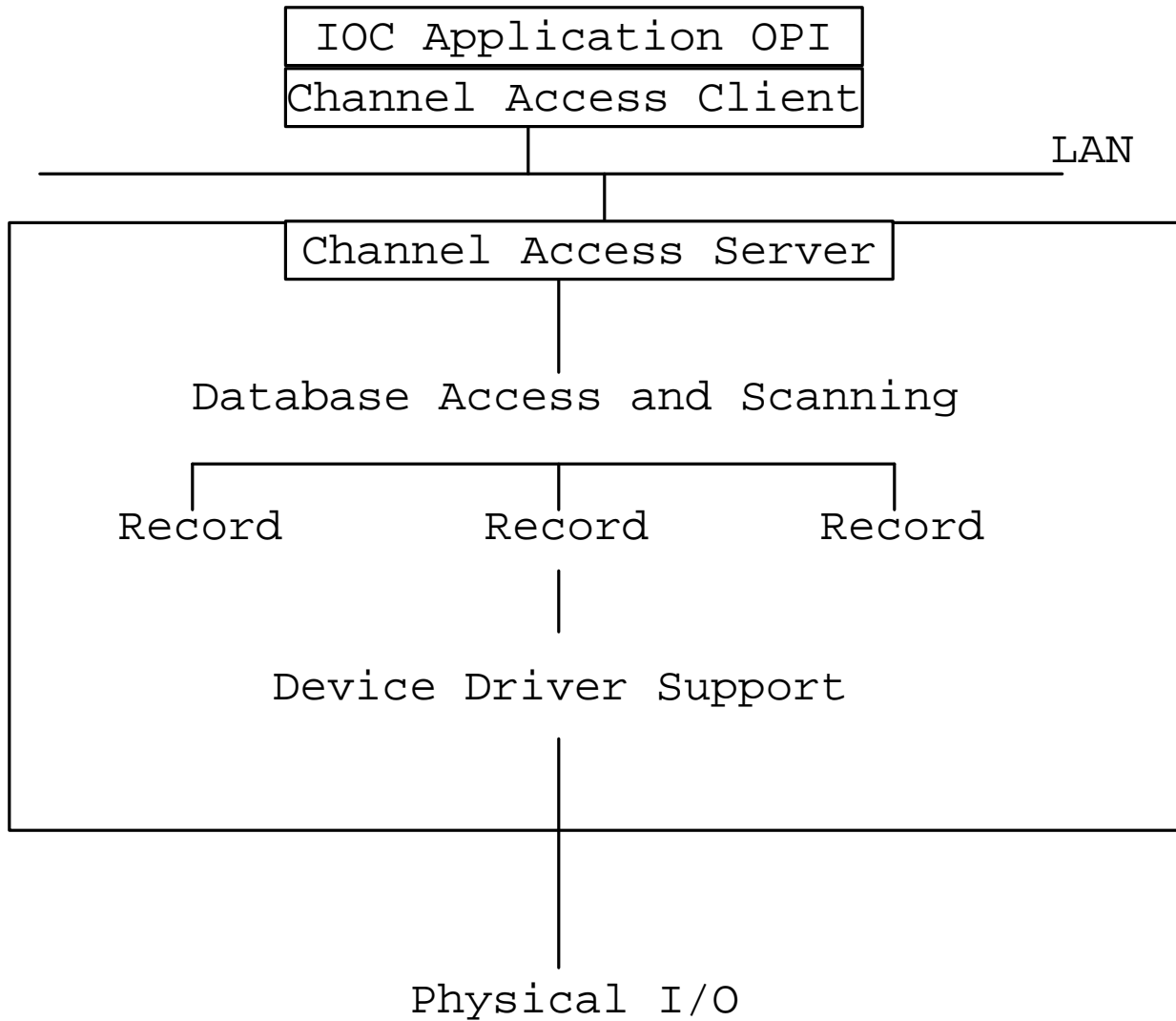
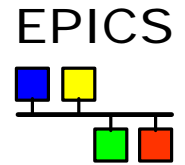
Flow of Control

Marty Kraimer
APS

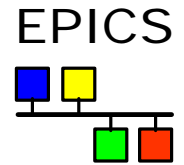
EPICS ARCHITECTURE PHYSICAL



EPICS ARCHITECTURE LOGICAL

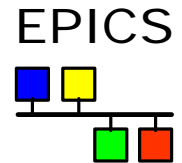


Summary



- ◆ IOC Initialization
 - ◆ Load various components
 - ◆ Connect to hardware
 - ◆ Start Channel Access Servers
- ◆ Record Processing
 - ◆ Hardware I/O monitor and control
 - ◆ Post events for CA clients
- ◆ Channel Access
 - ◆ Connect to channels
 - ◆ Get/Put requests
 - ◆ Monitor Requests
 - ◆ Connection Management

IOC Initialization

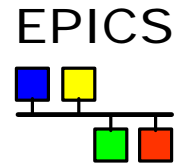


- ◆ `ld`
 - ◆ `iocCore`
 - ◆ `xxxApp` - device, driver, record support
- ◆ `dbLoadDatabase`

Load database definition files for menus, recordtypes, devices, and drivers
- ◆ `dbLoadRecords`, `dbLoadTemplates`

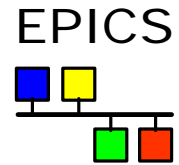
Load record instance files
- ◆ `iocInit`
 - ◆ Dynamically link to record, device, and driver support
 - ◆ Start general purpose tasks
log server, task watchdog, callback, dbCa, etc
 - ◆ Initialize drivers
 - ◆ Initialize records (two passes) and devices
 - ◆ Start scan tasks
 - ◆ Start access security
 - ◆ Start Channel Access Servers

Record Processing: Common Part



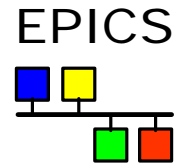
- ◆ Database Scanning calls dbProcess
 - ◆ Periodic
 - ◆ I/O Interrupt
 - ◆ Event
 - ◆ Passive
- ◆ Skip if already active
 - ◆ A lock set may have loops
 - ◆ The record may be asynchronous
- ◆ Skip if the record is disabled
- ◆ Call record support process routine
- ◆ Linked records (input, output, forward) that are process passive are processed recursively

Process: Synchronous I/O



- ◆ Set PACT true (processing active)
- ◆ Perform I/O operation
 - ◆ For hardware link call device support
 - ◆ For database link
 - ◆ linked records (except scalar NPP NMS input) are in a common lock set.
 - ◆ recursive processing may happen
 - ◆ inputs processed before data transfer
 - ◆ outputs processed after data transfer
 - ◆ For channel access link
 - ◆ inputs are monitored
 - ◆ outputs are handled via a separate task
- ◆ Check for record specific alarm conditions
- ◆ Call `db_post_event` for any fields that change because of record processing
- ◆ Request processing of forward link
- ◆ Set PACT false
- ◆ return

Process: Asynchronous I/O

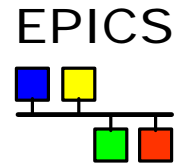


- ◆ **Asynchronous Start (Called by dbProcess)**
 - ◆ Set PACT True
 - ◆ Start I/O operation (normally via device support)
 - ◆ return leaving PACT true
- ◆ **Asynchronous Completion**

Record processing routine is called again

 - ◆ Complete I/O operation
 - ◆ Check for record specific alarm conditions
 - ◆ Call db_post_event for any fields that change because of record processing
 - ◆ Request processing of forward link
 - ◆ Set PACT false
 - ◆ return

Channel Access



◆ Search

- ◆ UDP broadcast PVNAME on local subnet (buffered)
- ◆ Each IOC has a CA broadcast listener
- ◆ IOC that has PV sends UDP message back to client
- ◆ TCP connection between each client/server
- ◆ For each client IOC establishes following tasks
 - ◆ Server task (all communication except monitors)
 - ◆ Event task which handles monitors, I.e. calls to `db_post_event`

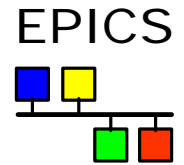
◆ Get

- ◆ Issue requests
- ◆ Wait for replies

◆ Put

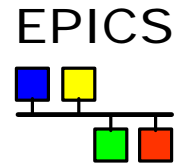
- ◆ Issue put. No response except for errors
- ◆ May cause record processing

Channel Access *continued*



- ◆ Add Event
 - ◆ Ask for monitor on PV (record.field)
 - ◆ Monitors happen because of calls to db_post_event
- ◆ Put Notify
 - ◆ Issue put and wait for completion
 - ◆ Completion when all records complete processing
- ◆ Connection Management
 - Automatic reconnect
- ◆ Access Security

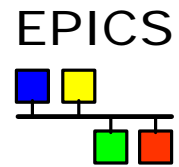
Example CA Put Notify



- ◆ Client issues put notify and flush
- ◆ IOC CA server accepts request
- ◆ dbPutNotify is called
- ◆ dbScanLock is called - Entire lock set is locked
- ◆ Value is written to record
- ◆ Record is processed (if passive and PP true in dbd file)
A set of records may start and some complete processing
- ◆ dbScanUnlock is called
- ◆ When all records complete processing the CA server is notified
- ◆ The CA server sends a message back to the client

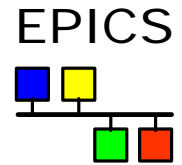
Example

Asynchronous Devices



- ◆ Serial and GPIB Devices
 - ◆ Are typically slow
 - ◆ Device support can NOT be synchronous
- ◆ Solution
 - ◆ Create separate task to manage requests
 - ◆ Device support is asynchronous
 - ◆ Async start queues a request to task
 - ◆ When task completes request it notifies device support which causes asynchronous completion.
 - ◆ Generic GPIB support is available
 - ◆ Several Serial support options are available
- ◆ Now lets look at CA Put Notify example
 - ◆ Database has series of linked records reading gpib/serial values.
 - ◆ rec1 -> rec2 -> rec3->...
 - ◆ What does CA Put Notify do?

CA put notify



- ◆ The following is sequence of events
 - ◆ Client issues put notify request
 - ◆ Server receives request
 - ◆ dbPutNotify is called
 - ◆ dbScanLock is called
 - ◆ Value is written to record rec1
 - ◆ rec1 starts processing
 - ◆ Device support queuesrequest to separate task
 - ◆ Record left with PACT true
 - ◆ dbscanUnlock is called
 - ◆ dbPutNotify waits
 - ◆ Sometime later
 - ◆ Task finishes request
 - ◆ rec1 has asynchronous completion
 - ◆ rec2 is processed and starts asynchronous phase
 - ◆ ... until all records in chain complete
 - ◆ dbPutNotify completes
 - ◆ CA sends message back to client